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	S. Da	PR	DA					
PR: Prevenble, DA: Destination Address, SA: Source Address, FCS: Frame Check Fig: Format of Ethernet CSMA/eD format Sequence.	PR:	and the second	A REAL PROPERTY	nation Ac	ldrun,	SA: Sou	rue Addres	<u>ل</u> ر ٥

(5 Q1) While difference between Token King and Token Bus. Token Bus Token King 1) The token is paned over the 1) The token is passed along Physical reing foremed by the the virtual ring of stations Stations and the Coaxial Cable Connected to a LAN. network. 215223 2) It is defined by IEEE 2) It is defined by IEEE 802.9 802.5 Standard. Standard 3) 9t, is not feasible to alculate 3) The maximum time for a token to reach a station the time fore taken transfere. can be calculated here A) The station are Connected 4) The underelying topology by ring topology, or that Connects the Stations Sometimes star topology. is eithere bus or tree topology station in LAN Station in LAN 2 5)usi 6053 (5) vaturilati mast be pretion Token Pamina Physical 400 reology 51:6 milio Direction 2 STO - Virtual Dirution 2 Ring + installing of token al faindo of Joken paning baden to hed asrobA Paning Fig: Token Ring, and as at Fig: Token Bus.

Q5) Write short motes on - (a) FODI (b) Taken Bus and Taken ving frame format (c) wirless LANI (d) Poll/Solect method ( controlled access method or polling)

Answers (a) Fiber Distributed Data Talasland (FDDI), 31 is a high performance fiber aptic taken ring. It is Similar to see is and IBM token ring, except the different that FDDI runs on fiber, not on Copper. FDDI Cable Consists of two fiber ring, one transmitting clockwise and othere transmitting Counter clockwise. If lither one breaks, the othere Can be used as a backup. if both breaks, the othere Can be used as a backup. if both break at the Same point, due to a fike or other accident in Cable duct, the two rings Can be yound into a Single ring. Each station Contains relays that Can be used to Joint the two rings or bypain the station in the event of Station problems.

134.00

Bytes 38 1 1 2 00 6 2 or 6 variable 4 1 1 B. Source Data ERC End France Addren M ERC delimiter studion 1 Bytes Preamble Shaft France Destimation delimiter Control Address Fig: Frame format of FDDI. (b) Taken Bus ken hing eframe format ?-1 1, 2016 6 2016 0-8182 Presentile Ethnet France Destimation Sources of the toy and the Bridden Bath CheckSum End delimiter delimiter 1 byte Fig: Token bus frame format.

Fjold Longth in bytes 1 1 1 6 6 20 4 1 bytes Shalf Access Prome Destination Securce Data Fes End Adverse Data Fes End delimiter.

delimiter Control delimiter

Taken

Fig: Taken King frame format.

(c) Divless lan : A wikeless local area onetwork (WLAN) is a wirless distribution method for two or more device that use high frequency readio waves and office often include an access point to the interenet. A WLAN allows useres to move around the Coverence area, often a home or small office, while maintaining a network Connection.

Advantages of WLAN

Taken wing from - format :

1) Flexibility. 1) Planning. Very Easy. 11) Design very Simple and portable. 11) Bobustness. transmission Technology. 12) Low Cost. for Designing. 13) Easy of use. 14) Easy of use. 15) Easy of use. 16) License free operation. (17) Low bandwidth.

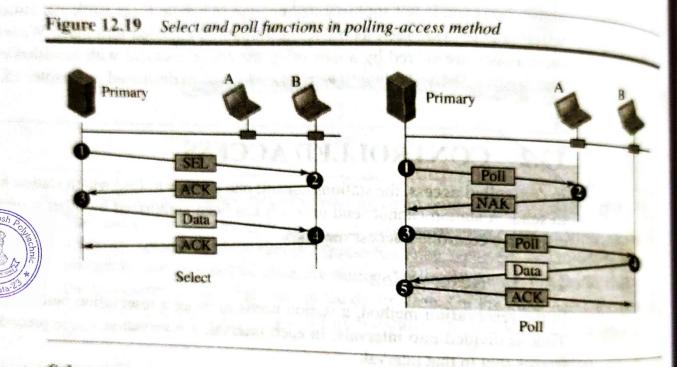


### DATA-LINA

# 5(d) 12.2.2 Polling (Controlled Access Metho)

(Polling works with topologies in which one device is designated as a primary station and (Polling works with topologies in which one device is designated as a primary station and (Polling works with topologies in which one device is designated as a primary station and (Polling works with topologies in which one device is designated as a primary station and (Polling works with topologies in which one device is designated as a primary station and (Polling works with topologies in which one device is designated as a primary station and (Polling works with topologies in which one device is designated as a primary station and (Polling works with topologies in which one device is designated as a primary station and (Polling works with topologies in which one device is designated as a primary station and (Polling works with topologies in which one device is designated as a primary station and (Polling works with topologies in which one device is designated as a primary station and (Polling works with topologies in which one device is designated as a primary station and (Polling works with topologies in which one device is designated as a primary station and (Polling works with topologies in which one device is designated as a primary station and (Polling works with topologies in which one device is designated as a primary station and (Polling works with topologies in which one device is designated as a primary station and (Polling works with topologies in which one device is designated as a primary station and (Polling works with topologies in which one device is device is designated as a primary station and (Polling works with topologies in which one device is device is designated as a primary station and (Polling works with topologies in which one device is device is device in which one device is device is device in which one device is device is device in which one de Polling works with topologies in which one the data exchanges must be made through the other devices are secondary stations. All data exchanges must be made through the ultimate destination is a secondary device. The the other devices are secondary stations. The destination is a secondary device. The primary device even when the ultimate destination is a secondary device. The primary devices follow its instructions) It is up to device the primary devices follow its instructions. primary device even when the utilitate devices follow its instructions) It is up to the primary device controls the link, the secondary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given the primary device is allowed to use the channel at a given th device controls the link, the secondary device is allowed to use the channel at a given time in many device to determine which device is allowed to use the channel at a given time in the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a session (see Figure 1) to the initiator of a primary device, therefore, is always the initiator of a session (see Figure 12.19). The method uses poll and select functions to prevent collisions. However, the drawback is the

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### Select

The select function is used whenever the primary device has something to send Remember that the primary controls the link. If the primary is neither sending and receiving data, it knows the link is available. If it has something to send, the primate device sends it. What it does not know, however, is whether the target device is pre pared to receive. So the primary must alert the secondary to the upcoming transmission and wait for an acknowledgment of the secondary's ready status. Before sending data the primary creates and transmits a select (SEL) frame, one field of which includes its address of the intended secondary.

## Pull

The poll function is used by the primary device to solicit transmissions from the solicit transmission of the solicit transmission ondary devices. When the primary device to solicit transmissions from device in turn if it has anything to receive data, it must ask (poil) call device in turn if it has anything to send. When the first secondary is approached responds either with a NAK frame if and. When the first secondary is approached a data frame) if it does. If the recommendation is approached to send or with data (in the form of the prime) a data frame) if it does. If the response is negative (a NAK frame), then the prime polls the next secondary in the response is negative (a NAK frame), then the public the response is positive (a data frame) until it finds one with data to send. the response is positive (a data frame), the primary reads the frame and returns a acknowledgment (ACK frame), verifying its receipt.